# Skoltech

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# SDN&NFV: Software-Defined Networks (SDN)

### Advanced Computer Networks

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# Part I: SDN



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## **Network problems**





Function Function

## **Operating System**

#### **Special hardware**

- Manufacturer dependent
- Errors in network protocol implementations
- Millions of Proprietary Closed Lines (6000+ RFC)
- High cost of equipment
- High cost of operation
- The difficulty of managing large networks
- Debugging complexity
- "Closed" equipment and software
- The difficulty of introducing new ideas
- Inefficiency in the use of hardware resources, energy inefficiency

## **Constantly increasing difficulty**



Source: http://www.telegeography.com/products/ip\_transit/index.php; http://www.ietf.org/

## **SDN** principles

## **S**oftware **D**efined **N**etworking (SDN)

- 1. Physically separate the network control plane from the data plane.
- 2. Logically centralized network management.
- 3. Open interface between the control plane and the data plane.
- New tools and features;
- Ease of administration;
- Openness to innovation and experimentation;
- IT market revolution

Creation	Platform Development	Implementations	Demonstrations	Market solutions	
2007 SDN Concept Stanford&Berkley)	2007-2010 Open Source Tools and Platforms: NOX&POX	2009-2012 Stanford GENI, Campuses Networks	2008-2012 SIGCOMM Open Networking Summit	2011-201? Manufacturers	
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## Forward to SDN



## **SDN** Architecture



## **SDN Advantages**



- Management flexibility
- Cheaper equipment (reduced CAPEX)
- Facilitate network management (OPEX reduction)
- Programmability, openness, innovation

## **SDN Use Case**



Decrease in power consumption in the data centers:

- Disabling Unused Switches and Channels Based on Collected Network Information
- ElasticTree (Stanford): up to 60% less energy
- Google application

## **Abstractions in IT**

Slowly developing, closed, expensive system. Small market



Specialized Programs

Specialized Operating System

**Specialized Hardware** 



## Abstractions B SDN?



## Part II: OpenFlow



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# OpenFlow



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## **OpenFlow Protocol**

Message Types:

#### • Constroller to Switch messages:

- Switch Configuration
- Management and status monitoring
- Stream Table Management
- Features, Configuration, Modify-State (flow-mod), Read-State (multipart request), Packet-out, Barrier, Role-Request

#### Symmetric Messages

- Sending messages in both directions
- Detecting Controller-Switch Connectivity Issues
- Hello, Echo
- Asymmetric Messages
  - Sending messages from the switch to the controller
  - Announce a change in network status, switch status
  - Packet-in, flow-removed, port-status, error

## **OpenFlow 1.0**

## **Flow Table**

Rı	ıle	Act	ion	Stats					
RL	ule	Act	ion	Packet and Byte Cour			ers		
RL	ule		Ecrwording pockets to port (a) with booder rewriting						
Rule		2. Forcing a packet to the controller							
Rule     3. Packet reset       4. Own extensions									
Switch Port	MAC src	MAC dst	Eth type	VLAN ID	IP Src	IP Dst	IP Prot	TCP sport	TCP dport
+ field ma	ask								

## **OpenFlow Rules**

#### Switching

Switch	MAC	MAC	Eth	VLAN	IP	IP	IP	TCP	TCP	Action
Port	src	dst	type	ID	Src	Dst	Prot	sport	dport	
*	*	00:1f:	*	*	*	*	*	*	*	port6

#### Flow Switching

Switch	MAC	MAC	Eth	VLAN	IP	IP	IP	TCP	TCP	Action
Port	src	dst	type	ID	Src	Dst	Prot	sport	dport	
port3	00:20	00:1f	0800	vlan1	1.2.3.4	5.6.7.8	4	17264	80	port6

#### Firewall

Switch Port	MAC src	2	MAC dst	Eth type	VLAN ID	IP Src	IP Dst	IP Prot	TCP sport	TCP dport	Action
*	*	*		*	*	*	*	*	*	22	drop

## **One or multiple tables**

# Table space explosion A, B, C, Y could be MAC or IP addresses, VLAN tags, MPLS labels etc. A A B C

#### Single-table abstraction may use table space inefficiently compared to multiple tables

#### OF 1.0 Single Table

A, Y1	
A, Y2	
A, Y3	
A, Y4	
A, Y5	
<u> </u>	
B, Y2	
B, Y3	
B, Y4	
-	
B, Y5	
C, Y1	
C, Y2	
C, Y3	
C, Y4	
O VE	
– C, 15	

## **OpenFlow 1.1**



- Moving packet only forward
- Transition: modifying a package, updating a set of actions, updating metadata

## **Group Table**



## Meter Table

- To implement QoS and speed limits
  - For each thread or group of threads
  - Keeps track of counter values
  - Actions: **drop** or **dscp** remark

Meter Identifier	Meter Bands	Counters
------------------	-------------	----------

Band Type	Rate	Counters	Type specific arguments
-----------	------	----------	-------------------------

# **Multiple Controllers**

- OpenFlow 1.2:
  - Multiple controllers
  - Controller Roles
- ROLES:
  - Roles: Master, Slave, Equal
  - By default: each controller in Equal role for each switch.
  - Role change:
     OFPT\_ROLE\_REQUEST
  - Role distribution



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## **OpenFlow Controller**

- Program, TCP / IP server, waiting for switch connections
- He is responsible for ensuring the interaction of the switch application.
- Provides important services (e.g. topology detection, host monitoring)
- The network OS API or controller provides the ability to create applications based on a centralized programming model.

# **OpenFlow Controllers List**

- Controller Implementations:
  - Nox, Pox, MUL, Ruy, Beacon, OpenDaylight, Floodlight, Maestro, McNettle, Flower, Runos
  - Different programming form Python to Haskell, Erlang
- For educations and experiments Pox.
- Developers Community
  - ONOS (Stanford)
  - OpenDayLight (Cisco)
- RUNOS Controller (ARCCN, Russia)
  - arccn.github.io/runos







## **Flow Installation Modes**

## **Reactive mode**



## **Flow Installation Modes**

## **Proactive mode**



## **Routing in SDN/OpenFlow Network**





#### Dynamic reconfiguration in the event of a network error.

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## **Topology Detection in SDN/OpenFlow**



## **OpenFlow Abstractions**

Control Plane – forwarding

Data Plane – match-action tables

# Part III: SDN/OpenFlow Use Cases

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## **SDN Use Cases**

Companies

Telecom operators and service providers

Data center and clouds

# **Enterprise Network**

- Modern companies have a complex network infrastructure:
  - A large number of network elements
  - Ramified topology
  - A set of different routing and security policies



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# Administration difficulties

- Network administrators are responsible for maintaining the network infrastructure:
  - Network engineers translate high-level policies into low-level teams
  - Manual configuration of all network devices
  - Limited Network Device Management Toolkit
  - Retraining for each vendor
- There are so-called SNMP management systems, but only monitoring of the status, and not management - the configuration is still in manual mode.

#### Router Management

```
1. Configure Static-routes/ACLs
2. Configure RIP
3. Configure OSPF
0. Exit
Select Menu Number [0-3]: 1
router> enable
router# configure terminal
router(config)# ip route 5.5.5.5 255.255.255.255 2.2.2.2
router(config)# write
Configuration saved...
router(config)# _
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```

# Goals

- Make the network manageable without manual access to equipment.
- 2. Improve network management abstraction.

# Web Interface

cience/projects/arccn/2015/ross15/deploy/enterprise.html



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## **WiFi Controller & SDN Networks**

#### The principle of interaction with SDN controller over Northbound API:


### WiFi Controller & SDN Networks

Test bed consists of:

- Three hardware OpenFlow switches Extreme Networks SummitX 460t
- Two TP-LINK access points •
- A laptop moving from one AP to another one and running ping command to the outside server •
- Both controllers run on the same server







- The legacy network needs in average 1.5 seconds to reconfigure, while the SDN/OpenFlow network doesn't bring additional delay.
- This is because the migration procedure in Chandelle requires less than 80ms and the OpenFlow controller has enough time to reconfigure the switches.
- Finally, we have more faster roaming with SDN/OpenFlow.

## Телеком

- 1. Intellectual Traffic Engineering:
- Choosing the best path in the graph
- Channel Failure Response
- Bandwidth Reservation



# Telecom

- How to put all this into practice?
  - Greenfield?
  - Traditional Network Integration Issues
    - You need to play along with the protocols of the traditional network, i.e. answer queries correctly.
    - The fewer joints with a traditional network, the better.
  - The problem of integration with existing management systems

### WAN segment (Service Provide)

#### Services:

- L2 transit for B2C, B2B/G (Internet, VPN)
- Fast backup path
- SLA
- IPTV multicast
- VolP
- Mobile backhaul

#### Before:

- IS-IS (RFC 1195)
- OSPF
- PIM-SSM
- PIM-SM
- LDP (RFC 3036)
- Targeted LDP
- BGP (PW для AToM)
- RSVP (RFC 2205)

- MPLS PW
- BGP (RFC 4271)
- MP-BGP (RFC 4760)
- MPLS-VPN (RFC 4364)
- MPLS (RFC 3031, 3032)

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#### After:

• L2 PW

application + stats (no encap)

- Bridge domain (no learning)
- Multicast (IGMP)
- L2 LAG, L3 ECMP
- H-QoS

# Datacenters/Clouds

- Increased utilization of equipment and channels
- Monitoring and stream optimization
- User Network Virtualization
- Load balancing
- Quality Assurance









## Datacenters/Clouds

- SDN:
  - Without OpenFlow, because they use OpenStack
    - ONLY virtual channels
    - Channels, tables, new VM, politics
  - With OpenFlow for managing of physical devices
    - Channel quality, bottleneck detection

#### **Network Virtualization**



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### Software Defined Data Center



## Part IV: OpenFlow Controllers

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## **OpenFlow Controller Architecture**



# **Requirements to SDN Controller**

### Performance

- Throughput
  - events per second
- Delay
  - us
- Reliability and security
  - 24/7

### Programmability

- Functionality: applications and services
- Programming interface

# List of OpenFlow Controllers(2013)

- NOX-Classic
- NOX
- Beacon
- Floodlight
- SNAC
- Ryu
- POX
- Maestro
- Trema

- Helios
- FlowER
- MUL
- McNettle
- NodeFlow
- Onix
- SOX
- Kandoo
- Jaxon

- Cisco ONE controller
- Nicira NVP Controller
- Big Network
   Controller
- IBM Programmable
   Network Controller
- HP SDN Controller
- NEC Programmable Flow

## **Experimental study**

- Performance
  - maximum number of processing requests
  - request processing time for a given load
- Scalability
  - change in performance indicators with an increase in the number of connections with switches and with an increase in the number of processor cores
- Reliability
  - the number of failures during testing at a given load profile
- Security
  - resistance to incorrectly formed OpenFlow protocol messages

# Comparison Results (2013)



- Maximum throughput 7.000.000 flows per second.
- Minimum delay time from 50 to 75 μs.
- Disadvantages:
  - Reliability of controllers raised questions
  - Performance was not enough (DC> 10M fps)

## **Productivity increase**

#### The most resource-intensive tasks:

- Interaction with OpenFlow switches:
  - multithreading usage;
  - thread loading accounting and rebalancing.
- Receiving OpenFlow packets from the channel:
  - reading packets from network card memory, bypassing the OS Linux network stack;
  - context switching;
  - virtual addresses.

## Performance (In-kernel Controller)



- Throughput: 30M fps
- Delay: **45us**

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# Programmability

- In the controller language [fast]
- In any language through the REST interface [slow]
- Special programming languages with another abstraction (e.g. Pyretic, Maple)

# NorthBound API

- NorthBound API the interface between the controller and applications
- Programming with OpenFlow is not an easy task!
  - It is difficult to perform independent tasks (routing, access control)
  - Low level abstraction
  - You need to remember the rules on the switches.
  - Unknown rule order for switches
- Application Portability Between Controllers



#### A.3.4.1 Modify Flow Entry Message

Modifications to a flow table from the controller are done with the OFPT\_FLOW\_MOD message:

```
/* Flow setup and teardown (controller -> datapath). */
struct ofp_flow_mod {
   struct ofp_header header;
   uint64_t cookie;
                                /* Opaque controller-issued identifier. */
   uint64_t cookie_mask;
                                 /* Mask used to restrict the cookie bits
                                    that must match when the command is
                                    OFPFC_MODIFY* or OFPFC_DELETE*. A value
                                    of 0 indicates no restriction. */
    /* Flow actions. */
    uint8_t table_id;
                                 /* ID of the table to put the flow in.
                                     For OFPFC_DELETE_* commands, OFPTT_ALL
                                     can also be used to delete matching
                                     flows from all tables. */
                                  /* One of OFPFC_*. */
    uint8_t command;
    uint16_t idle_timeout;
                                 /* Idle time before discarding (seconds). */
                                  /* Max time before discarding (seconds). */
    uint16_t hard_timeout;
                                  /* Priority level of flow entry. */
    uint16_t priority;
   uint32_t buffer_id;
                                  /* Buffered packet to apply to, or
                                     OFP_NO_BUFFER.
                                     Not meaningful for OFPFC_DELETE*. */
   uint32_t out_port;
                                  /* For OFPFC_DELETE* commands, require
                                     matching entries to include this as an
                                     output port. A value of OFPP_ANY
                                     indicates no restriction. */
   uint32_t out_group;
                                  /* For OFPFC_DELETE* commands, require
                                     matching entries to include this as an
                                     output group. A value of OFPG_ANY
                                     indicates no restriction. */
   uint16_t flags;
                                 /* One of OFPFF_*. */
   uint8_t pad[2];
    struct ofp_match match;
                                 /* Fields to match. Variable size. */
   //struct ofp_instruction instructions[0]; /* Instruction set */
}:
OFP_ASSERT(sizeof(struct ofp_flow_mod) == 56);
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```

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## **REST request example**

roote	pc-1: otal	~# cu	rl http Receive	://127 d % Xf	7.0.0 ferd	.1:8080 Averag	/wm/stat e Speed	icflower Time	trypusher Time	/list/0 Time	/00:00:00:24:e8:79:29:1a/json   json_pp -t dumpe e Current	er
100	670		670	0	0	Dload 69791	Upload	Total	Spent	Left	t Speed	
SVARI	- 1		670			09191			And the second second		1 14444	
18		100:00	0:00:24	:e8:79	9:29:1	la' =>	(					
							'drop-	flow' =>	• (			
									'prior	tity' =>	=> 32767,	
ties a									actic	ns' =>	> undef,	
									flags	· => 0,	0,	
									Vers:	on' =>	> _, _> _1	
									Imatel	t ab I		
											'dataLaverVirtualLanPrioritvCodePoint' => 0,	
											'wildcards' => 3145983,	
											'networkDestinationMaskLen' => 32,	
											'networkProtocol' => 0,	
											'transportSource' => 0,	
											'networkSourceMaskLen' => 32,	
											'dataLayerSource' => '00:00:00:00:00:00',	
											Instant SmaOfferminel => 0	
											'dataLaverDestination' => '00:00:00:00:00:00'	
											'inputPort' => 0.	
											'networkDestination' => '10.10.2.2',	
											'transportDestination' => 0,	
											'networkSource' => '10.10.1.2',	
											'dataLayerVirtualLan' => -1	
									20.038	1	},	
									COOKS	.e. =>	· 45035997289868789',	
									Lengt	ht -> 1	72	
									'out Po	set! =>	> -1.	
									'xid'	-> 0,		
									'type'	-> 'FI	FLOW MOD',	
									'hard'	'imeout'	t' → 0,	
									'idle1	'imeout'	t' => 0,	
									'comma	and' =>	> 0	
							<i>2</i> .					
							<u> </u>					
roote	pc-1:	-+ 1										

```
of13::FlowMod fm2; // Table 0: in port,VLAN=ar.ep.stag -> output(ar.ep.port)
fm2.table id(FORWARDING TABLE);
fm2.priority(100);
fm2.cookie(cookie);
fm2.xid(mgr->impl->xid);
fm2.buffer id(OFP NO BUFFER);
fm2.add oxm field(new of13::InPort(main route[0].port));
fm2.add oxm field(new of13::VLANVid(end switch list[0].ep list[0].stag | of13::OFPVID PRESENT));
of13::ApplyActions acts2;
acts2.add action(new of13::OutputAction(end switch list[0].ep list[0].port, 0));
fm2.add instruction(acts2);
mgr->get connection(sw1 dpid)->send(fm2);
of13::FlowMod fm3; // Table 0: VLAN=ar.ep.stag -> META=bbd id, GOTO 1
fm3.table id(FORWARDING TABLE);
fm3.priority(100);
fm3.cookie(cookie);
fm3.xid(mgr->impl->xid);
fm3.buffer id(OFP NO BUFFER);
fm3.add oxm field(new of13::VLANVid(end switch list[0].ep list[0].stag | of13::OFPVID PRESENT));
fm3.add instruction(new of13::WriteMetadata(domain id, 0xFFF));
```

```
fm3.add_instruction(new of13::GoToTable(LEARNING_TABLE));
```

```
mgr->get_connection(sw2_dpid)->send(fm3);
```

```
// Table 1: META=bbd_id, in_port=dr.ep0.port -> T0_CONTROLLER
for (auto ep : end_switch_list[1].ep_list) {
    of13::FlowMod fm4; // Table 1: META=bbd_id, in_port=dr.ep0.port -> T0_CONTROLLER
    fm4.table_id(LEARNING_TABLE);
    fm4.priority(100);
    fm4.cookie(cookie);
    fm4.xid(mgr->impl->xid);
    fm4.buffer_id(OFP_NO_BUFFER);
    fm4.add_oxm_field(new of13::Metadata(domain_id, 0xFFFF));
    fm4.add_oxm_field(new of13::InPort(ep.port));
    of13::ApplyActions acts4;
```

### Possible problems in OpenFlow controllers

Example of **the problem** with running several apps independently:

- Forwarding and Span apps. First app sends a flow over port 1, while second ones sends the same flow over port 5. Rules intersect with each other.
- Final rules order in the flow table is unknown.
- Packets will go using only the first rule. Thus, only one app will work. Conflict!
- We may to resolve such conflicts and some others. Just ip\_src:10.0.0.1 -> output:1,5!



## Part V: RUNOS Controller

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## **RUNOS – Network Operating System**

Network Management System - the first Russian SDN controller RUNOS

RUssian Network Operation System



#### **RUNOS Open version**

- on Github http://arccn.github.io/runos/
- Own base in C ++ 11/14, not Java
- goal: to simplify the development of network applications and not to forget about performance
- applications: topology, route, rebuilding in the event of a break, REST, WebUI, proactive loading of rules, redundancy Active-Passive







## **RUNOS – Network Operating System**

Network Management System - the first Russian SDN controller RUNOS

RUssian Network Operation System



- In-Kernel RUNOS version
  - Super performance 30 million events per second
  - Custom Application Development







## **RUNOS – Network Operating System**

Network Management System - the first Russian SDN controller RUNOS

RUssian Network Operation System

There are different controller options with a single database and a different set of services and applications

- Commercial RUNOS version for telecommunications operators
  - The base is the same as on Github. Customers can develop applications themselves. Learn from accessible materials
  - B2C, B2B services (p2p, mp2mp, multicast, etc.)
  - Active Standby Mode



**RUN@S** 



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## **RUNOS: Features**

- The problem of launching multiple applications, integration with applications of other developers
  - Static tuning of applications for themselves is required, the order and method of transferring information between them.
  - There is no mechanism for controlling and resolving conflicts between applications (generation of overlapping rules).
- In RUNOS, the task is to solve the above problems:
  - part of the configuration occurs automatically according to meta information, linking occurs dynamically
  - conflict resolution system developed
  - A wide range of services to simplify the development of new applications



#### Features:

space

- Algorithmic policies (rule generation)
- Client-friendly API using EDSL grammar (low level details are hidden inside the runtime – overloading, templates)
- Modules composition (parallel and sequential composition)

## **Release Descriptions**

#### Base:

- controller core
- topology building
- building a route through the entire network
- first version of the rule generation system
  - Priority allocation, combination of rules
  - LOAD, MATCH, READ abstraction
  - Based on maple
- Rest API (Floodlight compatible)
- WebUI (download monitoring, viewing tables, deleting and adding rules)
- Proactive loading rules
- Cold reservation
- ARP caching



## **Release Descriptions**

Version 0.6 is one of the last big releases.

- Full update of the controller core structure. There is no binding to a specific version of the OpenFlow protocol. Own model, expandable for any new fields, including those specific to equipment.
- Batch grammar for network applications. Simplifies the development of new applications.
  - "pkt[eth src] == eth addr"
  - "if (ethsrc == A || ethdst == B) doA else doB"
  - "test((eth\_src & "ff0.....0") == "....")"
  - "modify(ip\_dst >> "10.0.0.1")"
  - decision are "unicast()", "broadcast()", "drop()"
- Updating the rule generation system increased speed and improved rule generation (by the number of rules and the number of priorities).
- Test system.
- Runos-book detailed documentation and instructions for developing new applications.
- Applications: stp, arp, flow-manager

## **Open source RUNOS**

- Sources: <u>http://arccn.github.io/runos/</u>
  - Apache, version 2.0
- Tutorial s(Readme.md + Runos-book)
  - building, installation, running
  - First application tutorial
- Virtual Machine





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# New RUNOS Releases

- v0.6.1
  - Clean OpenFlow interface for programming switches (the ability to create rules yourself)
  - Updated REST: compatibility with Ryu, a library for Postman
- v0.7
  - Optimization of the rule generation system:
    - Global network vision
    - Optimization of work by number of FlowMod
    - New applications: corporate network
  - Improving the Web interface (transferring part of the functionality from the commercial version)

### Performance



- Throughput: 10 000 000 flows per sec
- Delay: **55 μs**



## Implementation

#### Keywords: C ++ 11/14/17, QT, Boost (asio, proto, graph)

The main third-party components:

- libfluid project (\_base, \_msg)
  - for interaction with switches and analysis of OpenFlow 1.3 messages
- libtins
  - parsing packets inside OpenFlow messages
- glog (google log)
  - multithreaded logging
- tcmalloc (google performance tools)
  - alternative faster implementation of malloc / free
- json11
  - parsing the configuration file
- boost graph
  - Topology storage, route search

## Parameters

```
Config (json):
"controller": {
 "threads": 4
},
"loader": {
 "threads": 3
},
"link discovery": {
 "poll-interval" : 10,
 "pin-to-thread" : 2
},
"learning switch": {
```

. . .

- The number of controller threads is set
  - for interacting with switches
  - for applications

### Application list

- their parameters (poll-interval)
- lock the thread of execution or select for exclusive use (pin-to-thread, ownthread)

## Architecture

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#### **Controller initialization:**

- Starting the desired number of 1. threads
- Launching Service Components 2.
- Launch applications and distribute 3. them by thread

Controller

Defining the order in which 4. applications process events

Workers



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## Part VI: Application Development for RUNOS Controller

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## First application – L2 learning



- L2 table where particularly host resides (host <->sw:port)
- A->B. What should we do on sw1?
  - Learn and broadcast
- B->A. What should we do on sw3?
  - Learn and unicast
- Advanced question: will it work for ping utilities? Ping 10.0.0.2 (assuming B has this IP)
  - Yes, arp (broadcast), ip (icmp) Advanced Computer Networks
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#### L2 learning table

Host	Switch:port
А	1:1
В	3:2

### Host Databases

```
class HostsDatabase {
   boost::shared_mutex mutex;
   std::unordered map<ethaddr, switch and port> db;
```

```
public:
```

```
void learn(uint64_t dpid, uint32_t in_port, ethaddr mac)
{
    LOG(INFO) << mac << " seen at " << dpid << ':' << in_port;
    {
        boost::unique_lock< boost::shared_mutex > lock(mutex);
        db[mac] = switch_and_port{dpid, in_port};
    }
}    .insert(...)
```

```
boost::optional<switch_and_port> query(ethaddr mac)
```

```
boost::shared_lock< boost::shared_mutex > lock(mutex);
auto it = db.find(mac);
if (it != db.end())
    return it->second;
else
    return boost::none;
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```

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## L2 forwarding application

```
// Get required fields
ethaddr dst mac = pkt.load(ofb eth dst);
db->learn(connection->dpid(),
          pkt.load(ofb in port),
          packet cast<TraceablePacket>(pkt).watch(ofb eth src));
auto target = db->query(dst mac);
// Forward
                                                     } else {
if (target) {
                                                         flow->broadcast();
    flow->idle timeout(60.0);
                                                         return PacketMissAction::Continue;
    flow->hard timeout(30 * 60.0);
    auto route = topology->computeRoute(connection->dpid(),
                                          target->dpid);
    if (route.size() > 0) {
        flow->unicast(route[0].port);
    } else {
        flow->idle timeout(0.0);
        LOG (WARNING) << "Path from " << connection->dpid()
            << " to " << target->dpid << " not found";</pre>
    return PacketMissAction :: Contindeanced Computer Networks
```

## Part VII: Distributed Control Plane

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# High Availability (HA)

- The network operates in 365/24/7 mode.
- The SDN control platform must be running continuously.
- The goal of high availability is to support the continued availability of the management platform and network applications.
- Causes of downtime: maintenance, software and hardware errors, hardware failure, attack, power outage, accident.

Availability koefficient, %	Downtime
99,999	5 min
99,99	52 min
99,9	8,7 h
99	3,7 days
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#### HA Active/Standby Techniques /1



## Restoration



#### **Features of the solution:**

- OpenFlow  $\geq 1.0$
- Corporate networks.
- Does not scale
- Incomplete utilization of computing resources

- Single controller failure
- Loss of connection between switch and controller
- Controller overload

### HA Active/Active Techniques /2



Active / Active Backup Strategies

- Asymmetric
- <u>Symmetrical</u>
- High complexity [Requires coordination of controllers, global state support]
- High Availability [Minimum Downtime]
- High utilization of computing resources

## Conclusion

- SDN is already actively used in the industry and is the main trend in the development of the telecom industry.
- SDN! = OpenFlow
  - SDN an approach to the separation of the data layer and the management level
  - OpenFlow is one of the implementations. Others, XMPP, SNMP, overlay.

"SDN means thinking differently about networking"

## Video about SDN

http://www.youtube.com/watch?v=GRVygzcXrM0